

## 9. SOME ANALYTIC RESULTS ABOUT PRIMES AND THE DIVISOR FUNCTION

**9.1. The Prime Number Theorem.** How frequent are the primes? At the end of the eighteenth century, Gauss and Legendre suggested giving up looking for a formula for the  $n$ th prime, and proposed instead estimating the number of primes up to  $x$ . So, define the prime-counting function  $\pi(x)$  by

$$\pi(x) = \sum_{\substack{p \leq x \\ p \text{ prime}}} 1.$$

Gauss conjectured on computational evidence that  $\pi(x) \sim \frac{x}{\log x}$ . This was proved by independently by Hadamard and de la Vallée Poussin in 1896, and became known as

**Theorem 9.1** (The Prime Number Theorem). *We have  $\pi(x) \sim \frac{x}{\log x}$  as  $x \rightarrow \infty$ .*

It turns out to be more convenient to work with

$$\theta(x) = \sum_{\substack{p \leq x \\ p \text{ prime}}} \log p,$$

which is called *Chebyshev's  $\theta$ -function*. In terms of this function it can be shown (not difficult) that the Prime Number Theorem is equivalent to the statement  $\theta(x) \sim x$  ( $x \rightarrow \infty$ ).

We won't prove PNT here, but instead a weaker version, and in terms of  $\theta(x)$ :

**Theorem 9.2.** *As  $x \rightarrow \infty$  we have*

$$(\log 2)x + o(x) < \theta(x) < (2 \log 2)x + o(x),$$

*so that*

$$0.6931x + o(x) < \theta(x) < 1.3863x + o(x).$$

## 9.2. Proof of Theorem 9.2.

### 9.2.1. The upper bound.

**Proposition 9.3.** *We have  $\theta(x) < (2 \log 2)x + O(\log^2 x)$ .*

*Proof.* Consider  $\binom{2n}{n}$ . By the Binomial Theorem, it is less than  $(1+1)^{2n} = 4^n$ . Also, it is divisible by all primes  $p$  with  $n < p \leq 2n$ , so

$$4^n > \binom{2n}{n} \geq \prod_{n < p \leq 2n} p = e^{\theta(2n) - \theta(n)}.$$

Hence  $\theta(2n) - \theta(n) \leq 2n \log 2$ .

Now if  $2n \leq x < 2n+2$  (i.e.,  $n \leq x/2 < n+1$ ) then  $\theta(x/2) = \theta(n)$  and

$$\theta(x) \leq \theta(2n) + \log(2n+1) \leq \theta(2n) + \log(x+1),$$

so that, for each  $x$ ,

$$\begin{aligned}\theta(x) - \theta(x/2) &\leq \theta(2n) + \log(x+1) - \theta(n) \\ &\leq 2n \log 2 + \log(x+1) \\ &\leq x \log 2 + \log(x+1).\end{aligned}$$

So (standard telescoping argument for  $x, x/2, x/2^2, \dots, x/2^k$  where  $x/2^{k-1} \geq 2$ ,  $x/2^k < 2$ ,  $\theta(x/2^k) = 0$ ):

$$\begin{aligned}\theta(x) &= \left(\theta(x) - \theta\left(\frac{x}{2}\right)\right) + \left(\theta\left(\frac{x}{2}\right) - \theta\left(\frac{x}{2^2}\right)\right) + \left(\theta\left(\frac{x}{2^2}\right) - \theta\left(\frac{x}{2^3}\right)\right) + \dots + \left(\theta\left(\frac{x}{2^{k-1}}\right) - \theta\left(\frac{x}{2^k}\right)\right) \\ &\leq \log 2 \left(x + \frac{x}{2} + \dots + \frac{x}{2^{k-1}}\right) + k \log(x+1) \\ &\leq 2x \log 2 + \lfloor \log_2 x \rfloor \log(x+1) \\ &\leq 2x \log 2 + O(\log^2 x).\end{aligned}$$

□

9.2.2. *The lower bound.* To obtain an inequality in the other direction, we look at

$$d_n = \text{lcm}(1, 2, \dots, n) = e^{\sum_{p^m \leq n} \log p}.$$

Define

$$\psi(x) = \sum_{\substack{p^m \leq x \\ p \text{ prime}}} \log p;$$

(i.e.,  $\log p$  to be counted  $m$  times if  $p^m$  is the highest power of  $p$  that is  $\leq x$ ). So  $d_n = e^{\psi(n)}$ .

**Lemma 9.4.** *We have  $\psi(x) < \theta(x) + 2x^{1/2} \log x + O(\log^2 x)$ .*

*Proof.* Now

$$\begin{aligned}\psi(x) &= \sum_{p \leq x} \log x + \sum_{p^2 \leq x} \log x + \sum_{p^3 \leq x} \log x + \dots \\ &= \theta(x) + \theta(x^{1/2}) + \theta(x^{1/3}) + \dots + \theta(x^{1/k}),\end{aligned}$$

where  $k$  is greatest such that  $x^{1/k} \geq 2$ , i.e.,  $k = \lfloor \log_2 x \rfloor$

$$\begin{aligned}&< \theta(x) + \log_2 x \theta(x^{1/2}) \\ &< \theta(x) + 2x^{1/2} \log x + O(\log^2 x),\end{aligned} \quad \text{using Prop. 9.3.}$$

□

Curious note: this  $k$  is the same one as in the proof of Prop. 9.3, though they have apparently different definitions.

We can now prove

**Proposition 9.5.** *We have  $\theta(x) \geq x \log 2 + O(x^{1/2} \log x)$ .*

*Proof.* Consider the polynomial  $p(t) = (t(1-t))^n$  on the interval  $[0, 1]$ . As  $t(1-t) \leq \frac{1}{4}$  on that interval (calculus!), we have

$$0 \leq p(t) \leq \frac{1}{4^n} \quad \text{on } [0, 1].$$

Writing  $p(t) = \sum_{k=0}^{2n} a_k t^k \in \mathbb{Z}[t]$ , then

$$\frac{1}{4^n} \geq \int_0^1 p(t) dt = \sum_{k=0}^{2n} \frac{a_k}{k+1} = \frac{N}{d_{2n+1}} \geq \frac{1}{d_{2n+1}},$$

for some  $N \in \mathbb{N}$ , on putting the fractions over a common denominator. Hence we have successively

$$\begin{aligned} d_{2n+1} &\geq 4^n \\ \psi(2n+1) &\geq 2n \log 2 && \text{on taking logs} \\ \theta(2n+1) &\geq 2n \log 2 - 2 \log(2n+1) \sqrt{2n+1} && \text{by Lemma 9.4} \\ \theta(x) &\geq x \log 2 + O(x^{1/2} \log x). \end{aligned}$$

□

Combining Propositions 9.3 and 9.5, we certainly obtain Theorem 9.2.

### 9.3. Some standard estimates.

**Lemma 9.6.** *For  $t > -1$  we have  $\log(1+t) \leq t$ , with equality iff  $t = 0$ .*

*For  $n \in \mathbb{N}$  we have  $n \log(1 + \frac{1}{n}) < 1$ .*

*Proof.* The first inequality comes from observing that the tangent  $y = t$  to the graph of  $y = \log(1+t)$  at  $t = 0$  lies above the graph, touching it only at  $t = 0$ . The second inequality comes from putting  $t = 1/n$  in the first inequality. □

**Lemma 9.7** (Weak Stirling Formula). *For  $n \in \mathbb{N}$  we have*

$$n \log n - n < \log(n!) \leq n \log n.$$

*Proof.* Now for  $j \geq 2$  we have

$$\begin{aligned} \log j &= j \log j - (j-1) \log(j-1) - (j-1) \log \left( 1 + \frac{1}{j-1} \right) \\ &= j \log j - (j-1) \log(j-1) - \delta_j, \end{aligned}$$

where  $0 < \delta_j < 1$ , using Lemma 9.6 for  $n = j - 1$ . So, on summing for  $j = 2, \dots, n$  we get

$$\begin{aligned} \log(n!) &= \sum_{j=2}^n \log j \\ &= \sum_{j=2}^n j \log j - (j-1) \log(j-1) - \delta_j \\ &= n \log n - \sum_{j=2}^n \delta_j \\ &= n \log n - \Delta, \end{aligned}$$

where  $0 < \Delta < n$ , since  $1 \log 1 = 0$  and all the other  $j \log j$  terms apart from  $n \log n$  telescope.  $\square$

**Proposition 9.8.** *We have*

$$\sum_{n \leq x} \frac{1}{n} = \log x + \gamma + O\left(\frac{1}{x}\right),$$

where  $\gamma = 0.577\dots$ , the Euler-Mascheroni constant.

*Proof.* Draw the graph of  $y = 1/t$  for  $t$  from  $0+$  to  $N+1$ , where  $N = \lfloor x \rfloor$ . On each interval  $[n, n+1]$  draw a rectangle of height  $1/n$ , so that these rectangles for  $n = 1, 2, \dots, N$  completely cover the area under the curve from  $t = 1$  to  $t = N+1$ . The pie-shaped pieces of the rectangles above the curve, when moved to the left to lie above the interval  $[0, 1]$ , are non-intersecting, and more than half-fill the  $1 \times 1$  square on that interval. Say their total area is  $\gamma_n$ . Then, as  $n \rightarrow \infty$ ,  $\gamma_n$  clearly tends to a limit  $\gamma$ , the Euler-Mascheroni constant.

The sum of the areas of the rectangles above  $[n, n+1]$  for  $n = 1, 2, \dots, N$  is clearly  $\sum_{n=1}^N 1/n$  (the total area of the parts of the rectangles below the curve). On the other hand, it is  $\int_1^{N+1} \frac{dx}{x} = \log(N+1)$  (the total area of the parts of the rectangles below the curve), plus  $\gamma_n$  (the total area of the parts of the rectangles above the curve). Hence

$$\sum_{n \leq x} \frac{1}{n} = \sum_{n=1}^N 1/n = \log(N+1) + \gamma_n.$$

Since  $\log(N+1) - \log x = O\left(\frac{1}{x}\right)$  and  $\gamma - \gamma_n = O\left(\frac{1}{x}\right)$  (check!), we have the result.  $\square$

**9.4. More estimates of sums of functions over primes.** Let us put  $\mathcal{P}_x = \prod_{p \leq x} \frac{1}{1-p^{-1}}$ . Then

**Proposition 9.9.** *We have  $\mathcal{P}_x > \log x$ .*

*Proof.* We have

$$\mathcal{P}_x = \prod_{p \leq x} \left(1 + \frac{1}{p} + \frac{1}{p^2} + \dots + \frac{1}{p^n} + \dots\right).$$

On multiplying these series together, we obtain a sum of terms that includes all fractions  $\frac{1}{n}$ , where  $n \leq x$ . This is simply because all prime factors of such  $n$  are at most  $x$ . Hence

$$\mathcal{P}_x > \sum_{n \leq x} \frac{1}{n} > \log x,$$

by Prop. 9.8. □

**Corollary 9.10.** *There are infinitely many primes.*

**Proposition 9.11.** *We have*

$$\sum_{p \leq x} \frac{1}{p} > \log \log x - 1.$$

*Proof.* We have

$$\begin{aligned} \log \mathcal{P}_x &= \sum_{p \leq x} \log \left( 1 + \frac{1}{p} + \frac{1}{p^2} + \frac{1}{p^3} + \cdots + \frac{1}{p^k} + \cdots \right) \\ &< \sum_{p \leq x} \frac{1}{p} + \sum_{p \leq x} \frac{1}{p(p-1)}, \end{aligned}$$

on applying Lemma 9.6 with  $t = \frac{1}{p} + \frac{1}{p^2} + \frac{1}{p^3} + \cdots + \frac{1}{p^k} + \cdots$ , and summing the GP, starting with the  $1/p^2$  term,

$$\begin{aligned} &< \sum_{p \leq x} \frac{1}{p} + \sum_{n=1}^{\infty} \frac{1}{(n+1)n} \\ &= \sum_{p \leq x} \frac{1}{p} + \sum_{n=1}^{\infty} \left( \frac{1}{n} - \frac{1}{n+1} \right) \\ &= \sum_{p \leq x} \frac{1}{p} + 1, \end{aligned}$$

because of the telescoping of  $\sum_{n=1}^{\infty} \left( \frac{1}{n} - \frac{1}{n+1} \right)$ . Hence

$$\sum_{p \leq x} \frac{1}{p} > \log \mathcal{P}_x - 1 > \log \log x - 1,$$

using Prop. 9.9. □

**Proposition 9.12.** *We have*

$$\sum_{p \leq x} \frac{\log p}{p} = \log x + O(1) \quad \text{as } x \rightarrow \infty.$$

*Proof.* Now from Problem Sheet 1, Q8, we have

$$n! = \prod_{p \leq n} p^{\left\lfloor \frac{n}{p} \right\rfloor + \left\lfloor \frac{n}{p^2} \right\rfloor + \dots},$$

so that (taking logs)

$$\begin{aligned} \log(n!) &= \sum_{p \leq n} \left( \left\lfloor \frac{n}{p} \right\rfloor + \left\lfloor \frac{n}{p^2} \right\rfloor + \left\lfloor \frac{n}{p^3} \right\rfloor + \dots \right) \log p \\ &= \sum_{p \leq n} \left\lfloor \frac{n}{p} \right\rfloor \log p + S_n, \end{aligned}$$

where

$$\begin{aligned} S_n &:= \sum_{p \leq n} \left( \left\lfloor \frac{n}{p^2} \right\rfloor + \left\lfloor \frac{n}{p^3} \right\rfloor + \dots \right) \log p \\ &\leq \sum_{p \leq n} \left( \frac{n}{p^2} + \frac{n}{p^3} + \dots \right) \log p \\ &= n \sum_{p \leq n} \frac{\log p}{p(p-1)} \\ &< n \sum_{k=1}^{\infty} \frac{\log(k+1)}{(k+1)k} \\ &= nc, \end{aligned}$$

for some positive constant  $c$ , since the last sum is convergent. Hence  $nc > S_n > 0$ . Also, for  $n = \lfloor x \rfloor$  we have

$$\begin{aligned} n \sum_{p \leq x} \frac{\log p}{p} &\geq \sum_{p \leq x} \left\lfloor \frac{n}{p} \right\rfloor \log p \\ &> \sum_{p \leq x} \left( \frac{n}{p} - 1 \right) \log p \\ &= n \sum_{p \leq x} \frac{\log p}{p} - \theta(x). \end{aligned}$$

Hence

$$n \sum_{p \leq x} \frac{\log p}{p} \geq \sum_{p \leq x} \left\lfloor \frac{n}{p} \right\rfloor \log p > n \sum_{p \leq x} \frac{\log p}{p} - O(x),$$

since  $\theta(x) = O(x)$ , by Theorem 9.2. Now add the inequality  $nc > S_n > 0$  to the above inequality, to obtain

$$n \sum_{p \leq x} \frac{\log p}{p} + nc > \log(n!) > n \sum_{p \leq x} \frac{\log p}{p} - O(x).$$

Dividing by  $n$ , and using the fact that  $\frac{\log(n!)}{n} = \log n - O(1)$  from Prop. 9.7, we have

$$\sum_{p \leq x} \frac{\log p}{p} + O(1) > \log n - O(1) > \sum_{p \leq x} \frac{\log p}{p} - O(1).$$

Hence

$$\sum_{p \leq x} \frac{\log p}{p} = \log x + O(1).$$

□

**9.5. The average size of the divisor function  $\tau(n)$ .** The following result is a way of saying that an integer  $n$  has  $\log n + 2\gamma - 1$  divisors, on average. Recall that  $\tau(n)$  is the number of (positive) divisors of  $n$ .

**Proposition 9.13.** *We have, as  $x \rightarrow \infty$ , that*

$$\sum_{n \leq x} \tau(n) = x \log x + (2\gamma - 1)x + O(\sqrt{x}).$$

*Proof.* Now

$$\begin{aligned} \sum_{n \leq x} \tau(n) &= \sum_{n \leq x} \sum_{\ell | n} 1 \\ &= \sum_{\ell \leq x} \sum_{\substack{n = k\ell \\ k \leq \frac{x}{\ell}}} 1 \\ &= \sum_{\ell \leq x} \left\lfloor \frac{x}{\ell} \right\rfloor, \end{aligned}$$

on recalling that  $\lfloor y \rfloor$  is the number of positive integers  $\leq y$ ,

$$\begin{aligned} &= 2 \sum_{\ell \leq \sqrt{x}} \left\lfloor \frac{x}{\ell} \right\rfloor - \lfloor \sqrt{x} \rfloor^2 && \text{by Q10, Problem Sheet 1} \\ &= 2 \sum_{\ell \leq \sqrt{x}} \frac{x}{\ell} - x + O(\sqrt{x}) \\ &= 2x \left( \log \sqrt{x} + \gamma + O\left(\frac{1}{\sqrt{x}}\right) \right) - x + O(\sqrt{x}) && \text{using Prop. 9.8} \\ &= x \log x + (2\gamma - 1)x + O(\sqrt{x}). \end{aligned}$$

□