A quadratic penalty algorithm for linear programming and its application to linearizations of quadratic assignment problems

### Julian Hall Ivet Galabova

School of Mathematics University of Edinburgh

Computational Optimization in Action

8 June 2018





# Solving LP problems: Crash start

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A \boldsymbol{x} \leq \boldsymbol{b}$   $\boldsymbol{x} \geq \boldsymbol{0}$ 

### Choosing the initial basis

- "Slack" basis is simple choice x<sub>0</sub>
- Standard crash aims for feasible vertex  $x_F$
- "Idiot" crash aims for near-optimal point  $\bar{x}^*$



# Solving LP problems: Crash start

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A \boldsymbol{x} \leq \boldsymbol{b}$   $\boldsymbol{x} \geq \boldsymbol{0}$ 

#### Choosing the initial basis

- "Slack" basis is simple choice x<sub>0</sub>
- Standard crash aims for feasible vertex  $x_F$
- "Idiot" crash aims for near-optimal point  $ar{x}^*$
- Idiot crash exists as code in clp
  - What is it?
  - (Why) does it work?
  - How good is it?



# Solving LP problems: Crash start

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A \boldsymbol{x} \leq \boldsymbol{b}$   $\boldsymbol{x} \geq \boldsymbol{0}$ 

#### Choosing the initial basis

- "Slack" basis is simple choice x<sub>0</sub>
- Standard crash aims for feasible vertex  $x_F$
- "Idiot" crash aims for near-optimal point  $ar{x}^*$
- Idiot crash exists as code in clp
  - What is it?
  - (Why) does it work?
  - How good is it?
- Google wanted to know!



- **Definition:** Forrest (2002)
  - Source code of clp

- **Definition:** Forrest (2002)
  - Source code of clp
- Dissemination: Forrest (2014)
  - "I gave a bad talk on it years ago"
  - "You minimize mu\*objective + sum of squared primal infeasibilities"
  - "This is done column by column... you just solve a quadratic to get new value"
  - "Periodically you reduce mu"

- **Definition:** Forrest (2002)
  - Source code of clp
- Dissemination: Forrest (2014)
  - "I gave a bad talk on it years ago"
  - "You minimize mu\*objective + sum of squared primal infeasibilities"
  - "This is done column by column... you just solve a quadratic to get new value"
  - "Periodically you reduce mu"
- Analysis: Forrest (2014)
  - "For many problems you finish with a small sum of infeasibilities and an objective a bit higher than the optimal one"

minimize f(x) subject to r(x) = 0

### Quadratic penalty method

- Minimize  $\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^T \mathbf{r}(\mathbf{x})$
- Decreasing sequence  $\{\mu^k\}$
- ${old x}^k o {old x}^*$  as  $k o \infty$
- Subproblems increasingly ill-conditioned as  $\mu^k$  decreases

minimize f(x) subject to r(x) = 0

### Quadratic penalty method

- Minimize  $\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^T \mathbf{r}(\mathbf{x})$
- Decreasing sequence  $\{\mu^k\}$
- ${old x}^k o {old x}^*$  as  $k o \infty$
- Subproblems increasingly ill-conditioned as  $\mu^k$  decreases

### Augmented Lagrangian method

Minimize

$$\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \boldsymbol{\lambda}^{\mathsf{T}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x})$$

- Decreasing sequence  $\{\mu^k\}$ •  $\lambda_i^{k+1} = \lambda_i^k + \mu^k \mathbf{r}(\mathbf{x}^k)$
- $\mathbf{x}^k o \mathbf{x}^*$  and  $\mathbf{\lambda}^k o \mathbf{\lambda}^*$  rapidly so ill-conditioning not an issue

minimize f(x) subject to r(x) = 0

### Quadratic penalty method

- Minimize  $\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^T \mathbf{r}(\mathbf{x})$
- Decreasing sequence  $\{\mu^k\}$
- $\pmb{x}^k 
  ightarrow \pmb{x}^*$  as  $k 
  ightarrow \infty$
- Subproblems increasingly ill-conditioned as  $\mu^k$  decreases

### Beale (1985)

- Quadratic form minimization as LP crash
- Implemented in SCICONIC

### Augmented Lagrangian method

### Minimize

$$\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \boldsymbol{\lambda}^{\mathsf{T}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x})$$

- Decreasing sequence  $\{\mu^k\}$ •  $\lambda_i^{k+1} = \lambda_i^k + \mu^k \mathbf{r}(\mathbf{x}^k)$
- $\mathbf{x}^k o \mathbf{x}^*$  and  $\mathbf{\lambda}^k o \mathbf{\lambda}^*$  rapidly so ill-conditioning not an issue

minimize f(x) subject to r(x) = 0

### Quadratic penalty method

- Minimize  $\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^T \mathbf{r}(\mathbf{x})$
- Decreasing sequence  $\{\mu^k\}$
- ${old x}^k o {old x}^*$  as  $k o \infty$
- Subproblems increasingly ill-conditioned as  $\mu^k$  decreases

### Beale (1985)

- Quadratic form minimization as LP crash
- Implemented in SCICONIC

### Augmented Lagrangian method

### Minimize

$$\phi(\mathbf{x},\mu) = f(\mathbf{x}) + \boldsymbol{\lambda}^{\mathsf{T}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x})$$

- Decreasing sequence  $\{\mu^k\}$ •  $\lambda_i^{k+1} = \lambda_i^k + \mu^k \mathbf{r}(\mathbf{x}^k)$
- $\mathbf{x}^k o \mathbf{x}^*$  and  $\mathbf{\lambda}^k o \mathbf{\lambda}^*$  rapidly so ill-conditioning not an issue

### Idiot algorithm

- Starts like augmented Lagrangian
- Finishes like quadratic penalty method

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A\boldsymbol{x} = \boldsymbol{b}$   $\boldsymbol{x} \ge \boldsymbol{0}$ 

### The Idiot algorithm: r(x) = Ax - b

Initialize 
$$\mathbf{x}^0 \ge \mathbf{0}$$
,  $\mu^1$ ,  $\lambda^1 = \mathbf{0}$   
For  $k = 1, 2, 3, ...K$   
 $\mathbf{x}^k = \arg\min_{\mathbf{x} \ge \mathbf{0}} h(\mathbf{x}) = \mathbf{c}^T \mathbf{x} + {\lambda^k}^T \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^k} \mathbf{r}(\mathbf{x})^T \mathbf{r}(\mathbf{x})$   
Possibly update  $\mu$ :  
 $\mu^{k+1} = \mu^k / \omega$ , for some factor  $\omega > 1$   
 $\lambda^{k+1} = \lambda^k$   
Else update  $\lambda$ :  
 $\mu^{k+1} = \mu^k$   
 $\lambda^{k+1} = \mu^k \mathbf{r}(\mathbf{x}^k)$   
End

# Idiot crash: Algorithm

• Solve subproblem

$$\min_{\mathbf{x} \ge \mathbf{0}} h(\mathbf{x}) = \mathbf{c}^{\mathsf{T}} \mathbf{x} + {\boldsymbol{\lambda}^{k}}^{\mathsf{T}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x}) \quad \text{where} \quad \mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$$

• Solve subproblem

$$\min_{\mathbf{x} \ge \mathbf{0}} h(\mathbf{x}) = \mathbf{c}^{\mathsf{T}} \mathbf{x} + \lambda^{k^{\mathsf{T}}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x}) \quad \text{where} \quad \mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$$

- Initially
  - Penalty parameter  $\mu^1$  ranges between 0.001 and 1000
  - Perform 20-30 "sample iterations", minimizing component-wise twice
  - Possibly abandon Idiot if 10% primal infeasibility reduction not achieved

• Solve subproblem

$$\min_{\mathbf{x} \ge \mathbf{0}} h(\mathbf{x}) = \mathbf{c}^{\mathsf{T}} \mathbf{x} + \lambda^{k^{\mathsf{T}}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x}) \quad \text{where} \quad \mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$$

- Initially
  - Penalty parameter  $\mu^1$  ranges between 0.001 and 1000
  - Perform 20-30 "sample iterations", minimizing component-wise twice
  - Possibly abandon Idiot if 10% primal infeasibility reduction not achieved
- Then, according to LP dimensions
  - Number of iterations K ranges between 30 and 200
  - $\mu$  is reduced (every 3 or 6 iterations) by  $\omega =$  0.333 (typically)

• Solve subproblem

$$\min_{\mathbf{x} \ge \mathbf{0}} h(\mathbf{x}) = \mathbf{c}^{\mathsf{T}} \mathbf{x} + {\boldsymbol{\lambda}^{k}}^{\mathsf{T}} \mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}} \mathbf{r}(\mathbf{x})^{\mathsf{T}} \mathbf{r}(\mathbf{x}) \quad \text{where} \quad \mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$$

- Initially
  - Penalty parameter  $\mu^1$  ranges between 0.001 and 1000
  - Perform 20-30 "sample iterations", minimizing component-wise twice
  - Possibly abandon Idiot if 10% primal infeasibility reduction not achieved
- Then, according to LP dimensions
  - Number of iterations K ranges between 30 and 200
  - $\mu$  is reduced (every 3 or 6 iterations) by  $\omega =$  0.333 (typically)
- Final value of  $\boldsymbol{\mu}$  is around machine precision
- $oldsymbol{\lambda}^k o oldsymbol{0}$  rapidly

## Idiot crash: How effective is it?

**Results:** Speed-up of the clp primal simplex solver when the ldiot crash is used and the percentage of solution time accounted for by the ldiot crash

# Idiot crash: How effective is it?

**Results:** Speed-up of the clp primal simplex solver when the ldiot crash is used and the percentage of solution time accounted for by the ldiot crash

	Best		Worst			
$Model^1$	Speed-up	ldiot (%)	$Model^1$	Speed-up	ldiot (%)	
Linf_520c	9.4	8.2	FOME12	1.1	0.1	
stp3d	6.5	0.9	KEN-18	1.0	0.7	
self	6.1	22.7	dfl001	1.0	0.1	
$\text{STORM}_{-}1000$	4.5	0.8	PDS-80	1.0	0.1	
nug15	4.2	0.1	maros-r7	0.9	7.8	
STORM-125	4.1	10.1	TRUSS	0.8	17.1	

# Idiot crash: How effective is it?

**Results:** Speed-up of the clp primal simplex solver when the ldiot crash is used and the percentage of solution time accounted for by the ldiot crash

	Best		Worst			
$Model^1$	Speed-up	ldiot (%)	$Model^1$	Speed-up	ldiot (%)	
Linf_520c	9.4	8.2	FOME12	1.1	0.1	
stp3d	6.5	0.9	KEN-18	1.0	0.7	
self	6.1	22.7	dfl001	1.0	0.1	
$\text{STORM}_{-}1000$	4.5	0.8	PDS-80	1.0	0.1	
nug15	4.2	0.1	maros-r7	0.9	7.8	
storm-125	4.1	10.1	TRUSS	0.8	17.1	

- Mean speed-up is 1.9; mean solution time accounted for by Idiot is 6%
- For only some problems does vanilla clp use the Idiot crash and primal simplex

[1: Results drawn from experiments on 30 Mittelmann benchmarking problems]

## Idiot crash: Effect on clp benchmark performance

**Results:** Performance of clp relative to cplex, gurobi and xpress Mittelmann (25/04/18)

Model	cplex	gurobi	xpress	clp
LINF_520C	495	574	255	35
NUG15	338	12	7	14
QAP12	26	1	1	5
QAP15	365	12	6	13
SELF	18	12	15	5

**Results:** Performance of clp relative to cplex, gurobi and xpress Mittelmann (25/04/18)

Model	cplex	gurobi	xpress	clp
Linf_520c	495	574	255	35
NUG15	338	12	7	14
QAP12	26	1	1	5
QAP15	365	12	6	13
SELF	18	12	15	5

- $\bullet~\mbox{For LINF}_520\rm{C},~\mbox{clp}$  is vastly faster
- For the three QAP linearizations, clp is very much faster than cplex
- For SELF, clp is significantly faster

# Idiot crash: Can it solve LPs?

### Results: Accuracy of final point after (up to) 200 Idiot iterations

- Residual  $||A\boldsymbol{x} \boldsymbol{b}||_2$
- Objective error  $\frac{|f-f^*|}{\max(1,|f^*|)}$

# Idiot crash: Can it solve LPs?

Results: Accuracy of final point after (up to) 200 Idiot iterations

- Residual  $||A\boldsymbol{x} \boldsymbol{b}||_2$
- Objective error  $\frac{|f-f^*|}{\max(1,|f^*|)}$

Best			Worst			
Model	Residual	Objective	Model	Residual	Objective	
NUG15	$2.1  imes 10^{-10}$	$3.7 \times 10^{-4}$	DBIC1	$3.8  imes 10^{-1}$	$8.5  imes 10^{-2}$	
MAROS-R7	$4.0  imes 10^{-9}$	$2.3 \times 10^{-5}$	STORM-125	$1.4\! imes\!10^{0}$	$1.2  imes 10^{-1}$	
PDS-100	$7.6  imes 10^{-10}$	$3.7 \times 10^{-4}$	TRUSS	$7.1\! imes\!10^{\!-1}$	$3.2 \times 10^{-1}$	
QAP15	$2.1  imes 10^{-10}$	$2.8 \times 10^{-3}$	MOD2	$3.9\! imes\!10^0$	$2.1  imes 10^{-1}$	
LP22	$1.1\! imes\!10^{-9}$	$1.3 \times 10^{-3}$	PILOT87	$2.1\! imes\!10^0$	$6.8  imes 10^{-1}$	
DFL001	$1.1 \!  imes \! 10^{-9}$	$3.7 \times 10^{-3}$	WORLD	$4.3\! imes\!10^0$	$5.5  imes 10^{-1}$	

# Idiot crash: Can it solve LPs?

Results: Accuracy of final point after (up to) 200 Idiot iterations

• Residual 
$$||A\boldsymbol{x} - \boldsymbol{b}||_2$$

• Objective error  $\frac{|f-f^*|}{\max(1,|f^*|)}$ 

Best			Worst			
Model	Residual	Objective	Model	Residual	Objective	
NUG15	$2.1  imes 10^{-10}$	$3.7 \times 10^{-4}$	DBIC1	$3.8  imes 10^{-1}$	$8.5  imes 10^{-2}$	
MAROS-R7	$4.0  imes 10^{-9}$	$2.3 \times 10^{-5}$	STORM-125	$1.4\! imes\!10^{0}$	$1.2  imes 10^{-1}$	
PDS-100	$7.6  imes 10^{-10}$	$3.7 \times 10^{-4}$	TRUSS	$7.1\! imes\!10^{\!-1}$	$3.2 \times 10^{-1}$	
QAP15	$2.1  imes 10^{-10}$	$2.8 \times 10^{-3}$	MOD2	$3.9\! imes\!10^0$	$2.1  imes 10^{-1}$	
LP22	$1.1\! imes\!10^{-9}$	$1.3 \times 10^{-3}$	PILOT87	$2.1\! imes\!10^0$	$6.8  imes 10^{-1}$	
DFL001	$1.1 \!  imes \! 10^{-9}$	$3.7 \times 10^{-3}$	WORLD	$4.3\! imes\!10^0$	$5.5  imes 10^{-1}$	

• Idiot crash clearly solves some problems to acceptable tolerances

• Objective error measure using  $f^*$  is not an optimality test

#### Accuracy measure



A convenient single quality measure for the point returned by the Idiot crash is

$$\mathsf{qual}(oldsymbol{x}) = \|Aoldsymbol{x} - oldsymbol{b}\| imes rac{|f-f^*|}{\mathsf{max}(1,|f^*|)}$$

No problems with low value of qual(x) have accurate optimal objective function value but large residual

### Accuracy and condition



- Clear correlation between accuracy of final point and cond(A)
- Quadratic assignment problems are particularly well conditioned

## Idiot crash: What can be proved?

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A\boldsymbol{x} = \boldsymbol{b}$   $\boldsymbol{x} \ge \boldsymbol{0}$ 

The Idiot objective is bounded below for bounded LP problems

The Idiot objective  $h^{k}(\mathbf{x}) = \mathbf{c}^{T}\mathbf{x} + \lambda^{k^{T}}\mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}}\mathbf{r}(\mathbf{x})^{T}\mathbf{r}(\mathbf{x})$ , where  $\mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$ , has positive semi-definite Hessian  $A^{T}A$ , but unboundedness of  $h^{k}(\mathbf{x})$  implies unboundedness of the LP

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A\boldsymbol{x} = \boldsymbol{b}$   $\boldsymbol{x} \ge \boldsymbol{0}$ 

The Idiot objective is bounded below for bounded LP problems

The Idiot objective  $h^{k}(\mathbf{x}) = \mathbf{c}^{T}\mathbf{x} + \lambda^{k}{}^{T}\mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}}\mathbf{r}(\mathbf{x})^{T}\mathbf{r}(\mathbf{x})$ , where  $\mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$ , has positive semi-definite Hessian  $A^{T}A$ , but unboundedness of  $h^{k}(\mathbf{x})$  implies unboundedness of the LP

#### The Idiot algorithm with exact solution of subproblems converges

**Theorem:** Suppose, that  $\mathbf{x}^k$  is the exact global minimizer of  $h^k(\mathbf{x})$  for each k = 1, 2... and that  $\{\mu^k\} \to 0$  as  $k \to \infty$ . Then every limit point of the sequence  $\{\mathbf{x}^k\}$  is a solution to the LP problem.

minimize 
$$f = \boldsymbol{c}^T \boldsymbol{x}$$
 subject to  $A\boldsymbol{x} = \boldsymbol{b}$   $\boldsymbol{x} \ge \boldsymbol{0}$ 

The Idiot objective is bounded below for bounded LP problems

The Idiot objective  $h^{k}(\mathbf{x}) = \mathbf{c}^{T}\mathbf{x} + \lambda^{k}{}^{T}\mathbf{r}(\mathbf{x}) + \frac{1}{2\mu^{k}}\mathbf{r}(\mathbf{x})^{T}\mathbf{r}(\mathbf{x})$ , where  $\mathbf{r}(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$ , has positive semi-definite Hessian  $A^{T}A$ , but unboundedness of  $h^{k}(\mathbf{x})$  implies unboundedness of the LP

#### The Idiot algorithm with exact solution of subproblems converges

**Theorem:** Suppose, that  $\mathbf{x}^k$  is the exact global minimizer of  $h^k(\mathbf{x})$  for each k = 1, 2... and that  $\{\mu^k\} \to 0$  as  $k \to \infty$ . Then every limit point of the sequence  $\{\mathbf{x}^k\}$  is a solution to the LP problem.

However: Subproblems are not (currently) solved exactly and ill-conditioning due to small  $\mu$  mitigates against it

## Idiot crash: What happens next?

- Final point is not a vertex (basic) solution
- clp performs
  - Crossover to get a basic solution
  - Primal simplex to get an optimal solution

# Idiot crash: What happens next?

- Final point is not a vertex (basic) solution
- clp performs
  - Crossover to get a basic solution
  - Primal simplex to get an optimal solution

Model	Speed-up	ldiot (%)	Residual	Objective
QAP15	4.0	0.1	$2.1  imes 10^{-10}$	$2.8 \times 10^{-3}$
NUG15	4.2	0.1	$2.1  imes 10^{-10}$	$3.7  imes 10^{-4}$
QAP12	2.5	0.6	$3.6  imes 10^{-10}$	$1.7\! imes\!10^{\!-1}$
$\text{STORM}_{1000}$	4.5	0.8	$5.9\! imes\!10^{-6}$	$5.9  imes 10^{-2}$
stp3d	6.5	0.9	$7.0  imes 10^{-5}$	$2.7  imes 10^{-2}$
PDS-100	2.5	5.4	$7.6  imes 10^{-10}$	$3.7 \times 10^{-4}$
$Linf_520c$	9.4	8.2	$1.1 \!  imes \! 10^{1}$	$9.1 \times 10^{-3}$

Idiot is a worthwhile crash, but relatively expensive to establish optimality!

• Know: If the Idiot crash yields a feasible point  $\bar{x}^*$  then

$$f^* \leq \boldsymbol{c}^T \bar{\boldsymbol{x}}^* = \bar{f}^*$$

• Know: If the Idiot crash yields a feasible point  $\bar{x}^*$  then

$$f^* \leq oldsymbol{c}^Toldsymbol{ar{x}}^* = oldsymbol{ar{f}}^*$$

• Consider: dual problem

maximize 
$$f_D = \boldsymbol{b}^T \boldsymbol{y}$$
 subject to  $A^T \boldsymbol{y} + \boldsymbol{s} = \boldsymbol{c}$   $\boldsymbol{s} \ge \boldsymbol{0}$ 

• Know: If the Idiot crash yields a feasible point  $\bar{x}^*$  then

$$f^* \leq oldsymbol{c}^{\, au} oldsymbol{ar{x}}^* = oldsymbol{ar{f}}^*$$

• Consider: dual problem

maximize 
$$f_D = \boldsymbol{b}^T \boldsymbol{y}$$
 subject to  $A^T \boldsymbol{y} + \boldsymbol{s} = \boldsymbol{c}$   $\boldsymbol{s} \ge \boldsymbol{0}$ 

• If the Idiot crash yields a feasible point  $ar{m{y}}^*$  then

$$ar{f}_D^* = oldsymbol{b}^Toldsymbol{ar{y}}^* \leq f^*$$

• Hence  $f^*$  lies in the interval  $[\bar{f}_D^*, \bar{f}^*]$ 

• Know: If the Idiot crash yields a feasible point  $\bar{x}^*$  then

$$f^* \leq oldsymbol{c}^{\, au} oldsymbol{ar{x}}^* = oldsymbol{ar{f}}^*$$

• Consider: dual problem

maximize 
$$f_D = \boldsymbol{b}^T \boldsymbol{y}$$
 subject to  $A^T \boldsymbol{y} + \boldsymbol{s} = \boldsymbol{c}$   $\boldsymbol{s} \ge \boldsymbol{0}$ 

• If the Idiot crash yields a feasible point  $ar{m{y}}^*$  then

$$ar{f}_D^* = oldsymbol{b}^Toldsymbol{ar{y}}^* \leq f^*$$

- Hence  $f^*$  lies in the interval  $[\bar{f}_D^*, \bar{f}^*]$
- Unfortunately: Values of  $\bar{f}_D^*$  don't (yet) have high accuracy

# Idiot crash: Application to quadratic assignment problem linearizations

### Quadratic assignment problem (QAP)

min 
$$f(X) = \sum_{i,j,k,l} a_{ik} b_{jl} x_{ij} x_{kl}$$
 s.t.  $X = [x_{ij}]_{n \times n} \in \Pi_n$ 

This is a MIQP problem with  $n^2$  binary variables and 2n constraints

# Idiot crash: Application to quadratic assignment problem linearizations

### Quadratic assignment problem (QAP)

min 
$$f(X) = \sum_{i,j,k,l} a_{ik} b_{jl} x_{ij} x_{kl}$$
 s.t.  $X = [x_{ij}]_{n \times n} \in \Pi_n$ 

This is a MIQP problem with  $n^2$  binary variables and 2n constraints

### QAP linearization (Adams and Johnson)

$$\begin{array}{ll} \min & f(X) = \sum_{i,j,\,k,\,l} a_{ik} b_{jl} y_{ijkl} \\ \text{s.t.} & \sum_{i} y_{ijkl} = x_{kl}, \ j, k, l = 1, \dots, n; \quad \sum_{j} y_{ijkl} = x_{kl}, \ i, k, l = 1, \dots, n \\ & y_{ijkl} \ge 0, \ i, j, k, l = 1, \dots, n; \quad X = [x_{ij}]_{n \times n} \in \Pi_n \end{array}$$

This is a MILP problem with  $n^2$  binary variables;  $n^4$  continuous variables  $y_{ijkl} = x_{ij}x_{kl}$ and  $n^4 + 2n^3 + 2n$  constraints.

# Idiot crash: Application to quadratic assignment problem linearizations

### Results: Performance after (up to) 200 Idiot iterations

Model	Rows	Columns	Optimum	Residual	Objective	Error	Time
NUG05	210	225	50.00	$9.4 \times 10^{-9}$	50.01	$1.5 \!  imes \! 10^{\!-4}$	0.04
NUG06	372	486	86.00	$7.8 \times 10^{-9}$	86.01	$1.2 \!  imes \! 10^{-4}$	0.11
NUG07	602	931	148.00	$7.9 \!  imes \! 10^{-9}$	148.64	$4.3 \times 10^{-3}$	0.25
NUG08	912	1613	203.50	$7.0  imes 10^{-9}$	204.41	$4.5  imes 10^{-3}$	0.47
NUG12	3192	8856	522.89	$8.8 \!  imes \! 10^{\! -9}$	523.86	$1.8\! imes\!10^{\!-\!3}$	2.58
NUG15	6330	22275	1041.00	$8.9\! imes\!10^{-9}$	1041.38	$3.7  imes 10^{-4}$	5.13
NUG20	15240	72600	2182.00	$7.5 \times 10^{-9}$	2183.03	$4.7 \times 10^{-4}$	14.94
NUG30	52260	379350	4805.00	$1.1 \!  imes \! 10^{-8}$	4811.41	$1.3 \times 10^{-3}$	82.28

- $\bullet~\mbox{Solution}$  of  $_{\rm NUG30}$  intractable using simplex or IPM on the same machine
- Idiot crash consistently yields near-optimal solutions
- Useful within a branch-and-bound solver?

## Idiot crash: Future work

#### Exact Idiot

- Convergence of component-wise search can be prohibitively slow
- Solve

$$\min_{\boldsymbol{x}\geq\boldsymbol{0}} h(\boldsymbol{x}) = \boldsymbol{c}^{\mathsf{T}}\boldsymbol{x} + \boldsymbol{\lambda}^{k^{\mathsf{T}}}\boldsymbol{r}(\boldsymbol{x}) + \frac{1}{2\mu^{k}}\boldsymbol{r}(\boldsymbol{x})^{\mathsf{T}}\boldsymbol{r}(\boldsymbol{x})$$

directly using conjugate gradient based approach

• Preconditioning?

## Idiot crash: Future work

### Exact Idiot

- Convergence of component-wise search can be prohibitively slow
- Solve

$$\min_{\boldsymbol{x}\geq\boldsymbol{0}} h(\boldsymbol{x}) = \boldsymbol{c}^{\mathsf{T}}\boldsymbol{x} + \boldsymbol{\lambda}^{k^{\mathsf{T}}}\boldsymbol{r}(\boldsymbol{x}) + \frac{1}{2\mu^{k}}\boldsymbol{r}(\boldsymbol{x})^{\mathsf{T}}\boldsymbol{r}(\boldsymbol{x})$$

directly using conjugate gradient based approach

• Preconditioning?

### Parallel Idiot

- Approximate or exact Idiot dominated by cost of forming Ax
- "Obvious" parallelism is memory bound
- Problem-specific code?

- Has been presented in algorithmic form for the first time
- Generally beneficial for the primal revised simplex method
- Converges to an optimal solution when subproblems are solved exactly
- Consistently and quickly yields near-optimal solutions of QAP linearizations intractable with simplex or IMP

- Has been presented in algorithmic form for the first time
- Generally beneficial for the primal revised simplex method
- Converges to an optimal solution when subproblems are solved exactly
- Consistently and quickly yields near-optimal solutions of QAP linearizations intractable with simplex or IMP

Slides: http://www.maths.ed.ac.uk/hall/COA18

Report: http://www.maths.ed.ac.uk/hall/GaHa18