The eventual image

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1. Overview

One-slide summary of talk

- Question: Is there a universal method for turning an *endomorphism* of an object into an *automorphism* of a (perhaps different) object?
- Answer: Yes, for trivial general reasons. There are both left and right universal methods.
- Surprise: When the objects concerned are 'finite', the left and right universal methods are the same.

More precise one-slide summary

Let $\ensuremath{\mathscr{C}}$ be a category.

Write $Endo(\mathscr{C})$ for the category in which:

- objects are pairs (X, T) with $X \in \mathscr{C}$ and $T: X \longrightarrow X$
- maps $(X, T) \longrightarrow (Y, S)$ are maps $f: X \longrightarrow Y$ such that Sf = fT.

Write $Auto(\mathscr{C})$ for the full subcategory of objects (X, T) where T is an automorphism.

For general reasons, $Auto(\mathscr{C}) \hookrightarrow Endo(\mathscr{C})$ usually has both adjoints.

The surprise: When the objects of $\mathscr C$ are 'finite', these adjoints are equal.

Examples

Let $\ensuremath{\mathscr{C}}$ be any of these categories:

- FinSet = finite sets
- FDVS = finite-dimensional vector spaces
- **CptMS** = compact metric spaces and distance-decreasing maps.

Then:

 $\begin{array}{rcl} \operatorname{Auto}(\mathscr{C}) \hookrightarrow \operatorname{Endo}(\mathscr{C}) \text{ has a simultaneous left and right adjoint,} \\ & & \operatorname{Endo}(\mathscr{C}) & \longrightarrow & \operatorname{Auto}(\mathscr{C}) \\ & & (X, T) & \longmapsto & (\operatornamewithlimits{im}^{\infty}(T), T'). \\ & & & \text{the eventual image of } T \end{array}$ Moreover, for all $(X, T) \in \operatorname{Endo}(\mathscr{C})$, the composite $& & (\operatorname{im}^{\infty}(T), T') \xrightarrow{\operatorname{counit}} (X, T) \xrightarrow{\operatorname{unit}} (\operatorname{im}^{\infty}(T), T') \end{array}$

is the identity.

${\mathscr C}$ has eventual images

Eventual images, explicitly

Let T be an endomorphism of an object X.

An eventual image of T is a retract

of (X, T) such that T' is an automorphism and:

- whenever S is an automorphism of an object Y, any map $(Y, S) \longrightarrow (X, T)$ factors uniquely through $(\operatorname{im}^{\infty}(T), T') \mapsto (X, T)$;
- whenever S is an automorphism of an object Y, any map $(X, T) \longrightarrow (Y, S)$ factors uniquely through $(X, T) \twoheadrightarrow (\operatorname{im}^{\infty}(T), T')$.

Composing gives an idempotent T^{∞} on X.

Plan for rest of talk

- 2. Sets
- 3. Vector spaces
- 4. Metric spaces
- 5. Unifying theorem

2. Sets

Sets

Let X be a finite set and $T: X \longrightarrow X$. Then

$$X\supseteq TX\supseteq T^2X\supseteq\cdots$$

Put
$$\operatorname{im}^{\infty}(T) = \bigcap_{n \in \mathbb{N}} T^n X = \bigcap_{n \in \mathbb{N}} \operatorname{im}(T^n).$$

Lemma

T restricts to an automorphism T' of $im^{\infty}(T)$.

This gives

$$(X, T)$$

 $(\operatorname{im}^{\infty}(T), T').$

Sets

Lemma

 $\{T, T^2, T^3, \ldots\}$ contains a unique idempotent, T^{∞} , say.

Lemma

 $\mathsf{im}(T^{\infty}) = \mathsf{im}^{\infty}(T).$

This gives maps

$$(X, T) \mathfrak{S} T^{\infty}$$

$$\bigwedge_{im^{\infty}(T), T')}$$

with the universal properties required.

3. Vector spaces

Vector spaces

Let X be a finite-dimensional vector space and $T: X \longrightarrow X$. Put

$$\operatorname{im}^{\infty}(T) = \bigcap_{n \in \mathbb{N}} \operatorname{im}(T^n), \quad \operatorname{ker}^{\infty}(T) = \bigcup_{n \in \mathbb{N}} \operatorname{ker}(T^n).$$

Lemma

T restricts to an automorphism T' of $im^{\infty}(T)$.

Lemma

$$X = \operatorname{im}^{\infty}(T) \oplus \operatorname{ker}^{\infty}(T).$$

This decomposition gives projection and inclusion maps

$$(X, T) \mathfrak{S} T^{\alpha}$$

$$\bigwedge_{i=1}^{i} \bigvee_{j=1}^{i}$$

$$(\mathsf{im}^{\infty}(T), T')$$

with the universal properties required.

(In fact, T^{∞} is a polynomial in T.)

4. Metric spaces

Metric spaces

Consider category of compact metric spaces and distance-decreasing maps $(d(f(x_1), f(x_2)) \leq d(x_1, x_2)).$

Background fact: For a map $T: X \longrightarrow X$ in this category,

$$T$$
 is distance-preserving $\iff T$ is invertible
 $\iff T$ is surjective.

Let $T: X \longrightarrow X$ be an endomorphism. Put $\operatorname{im}^{\infty}(T) = \bigcap_{n \in \mathbb{N}} \operatorname{im}(T^n)$. Lemma

T restricts to an automorphism T' of $im^{\infty}(T)$.

This gives

$$(X, T)$$

$$\bigwedge^{}_{\text{im}^{\infty}(T), T').}$$

Metric spaces

Lemma

The inclusion



has a canonical retraction.

The maps

$$(X, T) \mathfrak{S} T^{\infty}$$

$$\bigwedge_{m^{\infty}(T), T'}$$

have the universal properties required.

(In fact, T^{∞} is in the closure of $\{T, T^2, \ldots\}$, for a suitable topology.)

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5. Unifying theorem

The theorem

Let \mathscr{C} be a category with a factorization system. Call the left maps surjections (\rightarrow) and the right maps embeddings (\rightarrow), and suppose that:

• for an endomorphism T in \mathscr{C} ,

$$T$$
 is an embedding $\iff T$ is invertible
 $\iff T$ is a surjection;

limits of sequences ··· → · → · and colimits of sequences · → · → ··· exist and preserve factorizations.

Theorem

 \mathscr{C} has eventual images. In particular, $Auto(\mathscr{C}) \hookrightarrow Endo(\mathscr{C})$ has a simultaneous left and right adjoint.

Proof: Show that $\operatorname{im}^{\infty}(T)$ is limit *and* colimit of $\cdots \xrightarrow{T} X \xrightarrow{T} X \xrightarrow{T} \cdots$.

Examples of the theorem

C	embeddings	surjections
FinSet	injections	surjections
FDVS	injections	surjections
CptMS	distance-preserving maps	surjections
FDVS ^G (group rep'ns)	injections	surjections
FinGrp, FinRing, etc.	injections	surjections

Special properties of the eventual image

Proposition

Eventual image is tracelike: given maps

$$X \xrightarrow[g]{f} Y,$$

we have $\operatorname{im}^{\infty}(gf) \cong \operatorname{im}^{\infty}(fg)$.

Proposition

Eventual image is dynamical: given $T: X \longrightarrow X$, we have

$$\operatorname{im}^{\infty}(T) = \operatorname{im}^{\infty}(T^2) = \operatorname{im}^{\infty}(T^3) = \cdots$$

Open questions

- 1. Is there a more satisfactory general (enriched) setting?
- When is a left Kan extension equal to a right Kan extension?
 (The inclusion *i*: N → Z of additive monoids induces a functor [Z, C] → [N, C], which is just Auto(C) → Endo(C). If C has eventual images then left and right Kan extensions along *i* are equal.)
- What about discrete-time dynamical systems in less finite contexts?
 E.g. for the endomorphism z → z² of C ∪ {∞}, the eventual image should probably be {z ∈ C : |z| = 1} ∪ {0,∞}.