

Codensity monads

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These slides: [available on my web page](#)

Preview

This talk is about a canonical categorical construction: the **codensity monad** of a functor.

When applied to some familiar functors, it produces concepts including these:

- prime ideals and fields of fractions
- the radical of an integer
- ultrafilters and ultraproducts
- compact Hausdorff spaces and sheaves on them
- double dualization of vector spaces
- linearly compact vector spaces
- probability measures
- ...

The overall idea

Any functor with a left adjoint gives rise to a monad.

But even many functors *without* a left adjoint give rise to a monad.

That's the 'codensity monad' of the functor.

If the functor (G) does have a left adjoint (F) then the codensity monad is the familiar thing ($G \circ F$). But there are many other interesting and significant examples.

Plan

1. The definition, from many angles

2. Examples

1. The definition, from many angles

(Isbell, Ulmer; Applegate & Tierney, A. Kock)

The shape of the definition

Let $G: \mathcal{B} \longrightarrow \mathcal{A}$ be a functor.

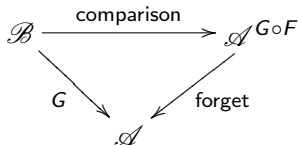
The **codensity monad** of G , if defined, is a monad T^G on \mathcal{A} .

(I'll define it later.)

It's not defined for *all* functors G , but it is defined for *many*, e.g. if \mathcal{B} is small and \mathcal{A} has small limits.

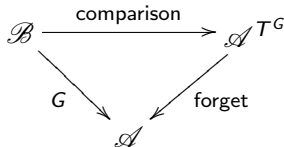
Motivation for the definition

Let $G: \mathcal{B} \longrightarrow \mathcal{A}$ be a functor that *does* have a left adjoint, F . We have categories and functors



and this is initial among all maps in \mathbf{CAT}/\mathcal{A} from G to a monadic functor.

Theorem (Dubuc) Let $G: \mathcal{B} \longrightarrow \mathcal{A}$ be a functor whose codensity monad T^G is defined. Then



is initial among all maps in \mathbf{CAT}/\mathcal{A} from G to a monadic functor.

Corollary Let G be a functor with a left adjoint, F . Then $T^G = G \circ F$.

The definition, via Kan extensions

Definition The **codensity monad** T^G of a functor G is the right Kan extension of G along itself (if it exists).

$$\begin{array}{ccc} \mathcal{B} & \xrightarrow{G} & \mathcal{A} \\ & \searrow G & \downarrow T^G \\ & & \mathcal{A} \end{array}$$

(A double arrow \Downarrow is shown between the top and bottom arrows.)

The monad structure on T^G comes from the universal property of Kan extensions.

E.g. the unit map $\eta: 1_{\mathcal{A}} \longrightarrow T^G$ is unique such that

$$\begin{array}{ccc} \mathcal{B} & \xrightarrow{G} & \mathcal{A} \\ & \searrow G & \downarrow T^G \\ & & \mathcal{A} \end{array} \quad \begin{array}{c} \Downarrow \eta \\ \eta \end{array} \quad 1 = \begin{array}{ccc} \mathcal{B} & \xrightarrow{G} & \mathcal{A} \\ & \searrow G & \downarrow \text{id} \\ & & \mathcal{A} \end{array} \quad \begin{array}{c} \Downarrow \text{id} \\ \text{id} \end{array} \quad 1$$

The definition, via ends

By the end formula for Kan extensions, the codensity monad T^G of $G: \mathcal{B} \longrightarrow \mathcal{A}$ is given by

$$T^G(A) = \int_B [\mathcal{A}(A, GB), GB] \in \mathcal{A}$$

($A \in \mathcal{A}$).

Here $[-, -]$ denotes a power: if $S \in \mathbf{Set}$ and $X \in \mathcal{A}$ then $[S, X] := \prod_{s \in S} X$.

The definition, via limits

Recall: $G: \mathcal{B} \longrightarrow \mathcal{A}$ and T^G is a monad on \mathcal{A} , with

$$T^G(A) = \int_B [\mathcal{A}(A, GB), GB].$$

Equivalently,

$$T^G(A) = \lim \left((A \downarrow G) \xrightarrow{\text{proj}} \mathcal{B} \xrightarrow{G} \mathcal{A} \right) = \lim_{\substack{B \in \mathcal{B}, \\ f: A \longrightarrow GB}} GB.$$

The definition, for **Set**-valued functors

If $\mathcal{A} = \mathbf{Set}$ then $G: \mathcal{B} \longrightarrow \mathbf{Set}$ and T^G is a monad on \mathbf{Set} , with

$$T^G(A) = \int_B [[A, GB], GB] = \lim_{\substack{B \in \mathcal{B}, \\ f: A \longrightarrow GB}} GB.$$

Equivalently,

$$T^G(A) = \{\text{natural transformations } G^A \longrightarrow G\}$$

where

$$\begin{array}{ccc} G^A: & \mathcal{B} & \longrightarrow \mathbf{Set}, \\ & B & \mapsto (GB)^A. \end{array}$$

E.g. Let $G: \mathbf{Grp} \longrightarrow \mathbf{Set}$ be the forgetful functor and let $A = n$ be a finite set. Then

$$\begin{aligned} T^G(n) &= \{\text{natural transformations } G^n \longrightarrow G\} \\ &= \{\text{maps } B^n \longrightarrow B \text{ defined for a generic group } B\} \\ &= \{n\text{-ary operations in the theory of groups}\} \\ &= \text{free group on } n \text{ generators.} \end{aligned}$$

2. Examples

First examples of codensity monads

- If G has a left adjoint F then $T^G = G \circ F$.
- Let $\mathcal{B} = \mathbf{1}$. A functor $\mathbf{1} \rightarrow \mathcal{A}$ is just an object X of \mathcal{A} . Its codensity monad T^X on \mathcal{A} is given by

$$T^X(A) = [\mathcal{A}(A, X), X].$$

This is the **endomorphism monad** $\mathbf{End}(X) = T^X$ of X .

It has the property that for *any* monad T on \mathcal{A} ,

$$T\text{-algebra structures on } X \leftrightarrow \text{monad maps } T \rightarrow \mathbf{End}(X).$$

(Compare group actions or representations.)

(Co)dense functors, 1

A functor $G: \mathcal{B} \longrightarrow \mathcal{A}$ is **codense** if its codensity monad is the identity.

This means that

$$A \cong \int_B [\mathcal{A}(A, GB), GB] = \lim_{f: A \longrightarrow GB} GB$$

naturally in $A \in \mathcal{A}$. Loosely,

‘every object of \mathcal{A} is a limit of objects in the image of G ’.

Often \mathcal{B} is a subcategory of \mathcal{A} and G is the inclusion. Then codensity means:

‘every object of \mathcal{A} is a limit of objects of \mathcal{B} ’.

Dually, \mathcal{B} is **dense** in \mathcal{A} if (loosely):

‘every object of \mathcal{A} is a colimit of objects of \mathcal{B} ’.

(Co)dense functors, 2

E.g.

- Let \mathbf{Mfd}_n be the category of smooth n -manifolds and smooth maps. Let \mathbf{Euc}_n be the subcategory of open subsets of \mathbb{R}^n and smooth open embeddings. Then $\mathbf{Euc}_n \hookrightarrow \mathbf{Mfd}_n$ is dense: ‘every manifold is a colimit of Euclidean patches’.
- Similarly, $\mathbf{CRing}^{\mathrm{op}} \cong (\text{affine schemes})$ is dense in the category of all schemes.

Warning: (Co)dense functors can behave counterintuitively! E.g. (Isbell):

- A composite of dense functors needn’t be dense.
- There’s an example of a full subcategory $\mathcal{B} \hookrightarrow \mathcal{A}$ such that \mathcal{B} is finite and \mathcal{A} is large (and not equivalent to any small category), but the inclusion is both dense and codense.

Rings and fields, 1

The inclusion $G: \mathbf{Field} \hookrightarrow \mathbf{CRing}$ has no left adjoint.

But it has a codensity monad! It's given by

$$T^G(A) = \prod_{\mathfrak{p} \in \mathrm{Spec}(A)} \mathrm{Frac}(A/\mathfrak{p})$$

($A \in \mathbf{CRing}$). Key ingredients of proof:

- For a homomorphism $\phi: A \rightarrow k$ into a field k , $\ker \phi$ is a prime ideal.
- Among all such homomorphisms with kernel \mathfrak{p} , the initial one is $A \rightarrow A/\mathfrak{p} \hookrightarrow \mathrm{Frac}(A/\mathfrak{p})$.

E.g.

$$\begin{aligned} T^G(\mathbb{Z}) &= \mathbb{Q} \times \mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/3\mathbb{Z} \times \mathbb{Z}/5\mathbb{Z} \times \cdots \\ T^G(\mathbb{Z}/n\mathbb{Z}) &= \mathbb{Z}/\mathrm{rad}(n)\mathbb{Z}, \end{aligned}$$

where $\mathrm{rad}(n)$ is the product of the distinct prime factors of n .

Rings and fields, 2

Some puzzles:

1. What are the algebras for T^G , the codensity monad of $\mathbf{Field} \hookrightarrow \mathbf{CRing}$?
2. What is $T^G(k[x])$, for a field k ?
3. What is the codensity monad of the forgetful functor $\mathbf{Field} \longrightarrow \mathbf{Set}$ (if it has one)?
4. What are *its* algebras?

Ultrafilters

Definition Let A be a set. An **ultrafilter** on A is a set \mathcal{U} of subsets of A such that whenever

$$A = A_1 \amalg \cdots \amalg A_n$$

($n \geq 0$), there is a unique i such that $A_i \in \mathcal{U}$.

Think of the elements of \mathcal{U} as the 'large' sets.

E.g.

- Given $a \in A$, get ultrafilter $\mathcal{U}_a = \{B \subseteq A : a \in B\}$.
- On a finite set A , every ultrafilter is of the form \mathcal{U}_a .
On an infinite set A , there are other ultrafilters—but proving this makes essential use of the axiom of choice.

Key properties of ultrafilters:

- they are upwards closed ($C \supseteq B \in \mathcal{U} \Rightarrow C \in \mathcal{U}$)
- they are closed under finite intersections
- if $B, B' \in \mathcal{U}$ then $B \cap B' \neq \emptyset$.

Ultrafilters from codensity monads

Theorem (Kennison and Gildenhuys) The codensity monad of the inclusion $G: \mathbf{FinSet} \hookrightarrow \mathbf{Set}$ is given by

$$T^G(A) = \{\text{ultrafilters on } A\}.$$

Sketch proof Recall that for $A \in \mathbf{Set}$,

$$\begin{aligned} T^G(A) &= \{\text{natural transformations } G^A \longrightarrow G\} \\ &= \{\text{natural families } (B^A \longrightarrow B)_{B \in \mathbf{FinSet}}\}. \end{aligned}$$

So an element of $T^G(A)$ is a machine that takes as input a finite set B and a function $A \longrightarrow B$, and produces as output an element of B .

Equivalently, it's a way of selecting, for each partition of A into finitely many pieces, one of those pieces.

Equivalently, it's an ultrafilter on A .

Compact Hausdorff spaces from codensity monads

We just saw that $A \mapsto \{\text{ultrafilters on } A\}$ is a monad on **Set**.

Theorem (Manes) The algebras for the ultrafilter monad are the compact Hausdorff spaces.

Background Let X be a topological space. Ultrafilters on X can be viewed as 'generalized sequences'.

An ultrafilter \mathcal{U} on X **converges** to $x \in X$ if every neighbourhood of x belongs to \mathcal{U} .

Facts:

- X is compact \iff every ultrafilter on X converges to *at least* one point.
- X is Hausdorff \iff every ultrafilter on X converges to *at most* one point.
- X is compact Hausdorff \iff every ultrafilter on X converges to *exactly* one point.

So if X is compact Hausdorff, get function $\lim: T^G(X) \longrightarrow X$.

Double duals

We saw that the codensity monad of $\mathbf{FinSet} \hookrightarrow \mathbf{Set}$ was something interesting.

What about $\mathbf{FVect} \hookrightarrow \mathbf{Vect}$?

Theorem

- The codensity monad of $\mathbf{FVect} \hookrightarrow \mathbf{Vect}$ is double dualization $(\)^{**}$.
- The algebras of $(\)^{**}$ are the **linearly compact vector spaces** (certain topological vector spaces).

Details omitted! This is the linear analogue of the story of ultrafilters and compact Hausdorff spaces.

Probability measures

Let **Msbl** be the category of measurable spaces (sets equipped with a σ -algebra).

Let \mathcal{B} be the subcategory consisting of the single object

$$B = \{\text{sequences in } [0, 1] \text{ converging to } 0\}$$

and the affine maps $B \longrightarrow B$ (those preserving convex combinations).

Theorem (Avery) The codensity monad of $\mathcal{B} \hookrightarrow \mathbf{Msbl}$ is the Giry monad,

$$T^G(A) = \{\text{probability measures on } A\}.$$

Moral

Whenever you meet a functor,
ask
“What is its codensity monad?”

References

The paper titles below are clickable links.

- Tom Leinster, [Codensity and the ultrafilter monad](#), *Theory and Applications of Categories* 28 (2013), 332–370.
- Tom Avery, [Codensity and the Giry monad](#), *Journal of Pure and Applied Algebra* 220 (2016), 1229–1251.

Other references, including references to work attributed to other authors in these slides, can be found in ‘Codensity and the ultrafilter monad’.